

# Solution to Exam in TSIN02 Internetworking

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1. See lecture 3, starting at slide 6.
2.  $\log_2(131072) = 17$ . 173.75.17.30=10101101.01001011.00010001.00011110. The closest subnet to this address which can contain 131072 addresses is thus 10101101.01001100.00000000.00000000-10101101.01001101.11111111.11111111, i.e., 173.76.0.0-173.77.255.255/15.  
This subnet is then subdivided into 10101101.01001100.00000000.00000000-10101101.01001100.11111111.11111111, and 10101101.01001101.00000000.00000000-10101101.01001101.11111111.11111111, i.e., 173.76.0.0-173.76.255.255/16, and 173.77.0.0-173.77.255.255/16.
3. Following the whiteboard notes of lecture 5, or Section 11.2.3 of 20Q, we have that
  - a) 25 SEK
  - b) -200 SEK
  - c) Under flat rate, the customers are unhappy since the net utility is negative. On the other hand, with usage-based pricing, the users are happy since the net utility is positive. This follows from that the underlying utility function is based on a user poll where the users tell how much they are willing to pay for different amounts of data.
4.
  - a)  $m = 3$ .
  - b)  $r = 500$ .
  - c)  $m = 3 \geq 2n - 1 = 2 \cdot 2 - 1 = 3$ , so the network is non-blocking.
  - d) We want to maintain the non-blocking property in (c). At the same time, we want to reduce the number of switch connections of the  $500 \times 500$ -middle switches. One choice is a  $(n,m,r)=(2,3,250)$ -Clos network.

See lecture 7 for more details!

5. a) The optimization problem is as follows:

$$\begin{aligned} & \underset{\mathbf{x}}{\text{maximize}} && F_{\beta,\lambda}(\mathbf{x}) = f_{\beta} \cdot \left(\sum_i x_i\right)^{1/\lambda} \\ & \text{subject to} && x_1 + x_4 \leq 1, \\ & && x_2 + x_4 \leq 1, \\ & && x_3 + x_5 \leq 1, \\ & && x_i \geq 0, \quad i = 1, \dots, 5. \end{aligned} \tag{1}$$

b) When  $1/\lambda = 0$ , the objective function is the fairness function, and equal allocation obtains the optimal value. When  $\frac{1}{\lambda}$  is increased, the weighting of the efficiency component increases and throughput is prioritized instead of fairness. The long session  $x_4$  is penalized and short sessions  $x_1$  and  $x_2$  are given priority. Sessions  $x_3$  and  $x_5$  have the same length, so equal allocation is both fair and efficient and hence optimal for all values of  $1/\lambda$ . The sum throughput is 2.5 for  $\frac{1}{\lambda} = 0$ , and 3 for  $\frac{1}{\lambda} = 10$ , which supports our interpretation above.

6. We use the collection of formulas to get

$$\text{MSE} = \frac{1 - a^{2|k|}}{1 - a^2} \quad (2)$$

and with  $k = 1$ , we have  $\text{MSE}=1$  with interleaving.

Without interleaving, we have very large  $k$  for most pixels. Hence, most pixels will have a distortion

$$\text{MSE} \approx \frac{1}{1 - a^2}, \quad (3)$$

and so the  $\text{MSE} \approx 2$  without interleaving.